

# Conflict-Based Diagnosis of Discrete-Event Systems

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We define a conflict-based diagnosis theory for discrete event systems

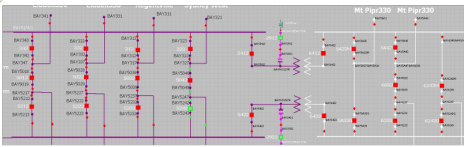
- Compatible with the existing conflict-based diagnosis for circuits (Reiter theory)
- Efficient (solve many unsolved problems)
- Applicable to more frameworks (e.g. hybrid systems)

- 1 Example
- 2 Diagnosis
- 3 Consistency-Based Diagnosis
- 4 Validation

## TransGrid Network



● 10k components



## Alarm Log (extract)

```

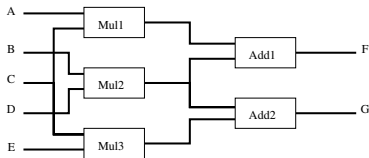
Date System_Time Event Text
2/07/2009 10:47:27 BAYSWTR PS 023 NO4 GEN UNIT STATUS OFF
2/07/2009 10:47:27 BAYSWTR330 330 SYD WEST 322 CB --OPENED--
2/07/2009 10:47:27 BAYSWTR330 330 NO4 BY/CUP 5042 CB --OPENED--
2/07/2009 10:47:27 BAYSWTR330 330 NO4 GEN TX 5242 CB --OPENED--
2/07/2009 10:47:27 BAYSWTR330 CONTROL SYSTEM LAN FAULT ALARM
2/07/2009 10:47:27 BAYSWTR PS 023 NO4 GEN 2242 CB --OPENED--
2/07/2009 10:47:28 LIDDELL330 330 BAYSWTR330 332 CB --OPENED--
2/07/2009 10:47:28 LIDDELL330 330 BAYSWTR330 342 CB --OPENED--
2/07/2009 10:47:28 LIDDELL330 330 NO2 BY/CUP 5022 CB --OPENED--
2/07/2009 10:47:28 LIDDELL330 330 NO3 BY/CUP 5032 CB --OPENED--
2/07/2009 10:47:28 WANG330 FAULT RECORDER OPERATED ALARM
2/07/2009 10:47:28 BAYSWTR330 330 MAIN BUS BAR KV Limit 5 Low
2/07/2009 10:47:28 BAYSWTR330 330 GEN BUS BAR KV Limit 5 Low
2/07/2009 10:47:28 WANG330 BU SUBSTATION MISC EQUIPMENT FAIL ALARM
2/07/2009 10:47:28 SYD WEST 330 BAYSWTR330 322B B CB --OPENED--
2/07/2009 10:47:28 SYD WEST 330 BAYSWTR330 322A A CB --OPENED--
2/07/2009 10:47:28 MT PIPR330 330 FAULT RECORDER OPERATED ALARM
2/07/2009 10:47:28 ERARING500 SUBSTATION MISC EQUIP FAIL ALARM
2/07/2009 10:47:28 MT PIPR330 500 B BUS BAR KV Limit 3 Low
2/07/2009 10:47:28 BAYSWTR330 330 NO3 BY/CUP 5032 CB --OPENED--
2/07/2009 10:47:28 BAYSWTR330 330 NO3 GEN TX 5232 CB --OPENED--
2/07/2009 10:47:28 BAYSWTR330 330 REGENTVILE 312 CB --OPENED--
2/07/2009 10:47:28 BAYSWTR PS 023 NO3 GEN 2232 CB --OPENED--

```

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# Model-Based Diagnosis

## Static Systems



$$A = B = E = 3$$

$$F = 10$$

$$C = D = 2$$

$$G = 12$$

- **Model** Formula  $\Phi_M$  involving *Ab* literals
- **Observation** Formula  $\Phi_O$
- **Possible behaviours**  $\Phi_M \wedge \Phi_O$
- **Diagnosis** Projection on the *Ab* literals:  $\exists X. \Phi_M \wedge \Phi_O$   
where  $X$  are the non *Ab* literals, rewritten in prime implicants

$$Ab(Mul1) \vee Ab(Add1) \vee (Ab(Mul2) \wedge Ab(Mul3)) \\ \vee (Ab(Mul2) \wedge Ab(Add2))$$

AUTOMATON

SEQUENCE OF  
OBSERVATIONS

- **Model** Language  $\mathcal{L}_M$  involving  $\Sigma_f$  events
- **Observation** Language  $\mathcal{L}_O$  involving only observable events  $\Sigma_O$
- **Possible behaviours**  $\mathcal{L}_M \cap \mathcal{L}_O$
- **Diagnosis** Projection on the  $\Sigma_f$  events and minimisation (removes non minimal words)

$$\mathcal{L}_\Delta = \text{Minimisation}(\text{Proj}_{\Sigma_f}(\mathcal{L}_M \cap \mathcal{L}_O))$$



### Static Systems

- **Model** Formula  $\Phi_M$
- **Observation** Formula  $\Phi_O$
- **Possible behaviours**  
 $\Phi_M \wedge \Phi_O$
- **Diagnosis** Projection on the *Ab* literal + prime implicants

### Discrete Event Systems

- **Model** Language  $\mathcal{L}_M$
- **Observation** Language  $\mathcal{L}_O$
- **Possible behaviours**  
 $\mathcal{L}_M \cap \mathcal{L}_O$
- **Diagnosis** Projection on the  $\Sigma_f$  events and minimisation



Boum!

## Static Systems

The size of the formula is exponential in the number of state variables

→ Compilation Map (Darwiche et al.), BDD, sd-DNNF, Cone-based diagnoser, etc.

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# Boum!

## DES

The size of the automata is exponential in the number of components

→ Decentralised / Distributed approach, Junction Trees, Specialised diagnosers, etc.

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Check carefully-chosen hypotheses until the diagnosis is found

- We do not compute all diagnosis candidates
- We compute only one representative of each candidate
- For each test, we derive useful information from the hypothesis at hand

## Static Systems

- $\Phi_h$  is a conjunct defined on all *Ab* literals
- $h$  is a candidate iff

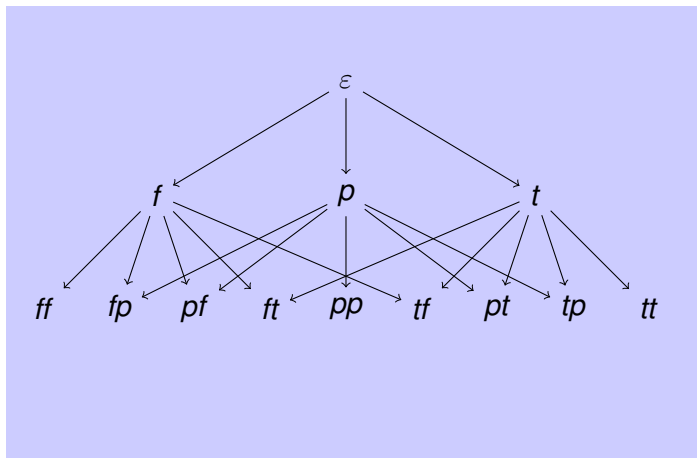
$$\Phi_M, \Phi_O, \Phi_h \not\models \perp$$

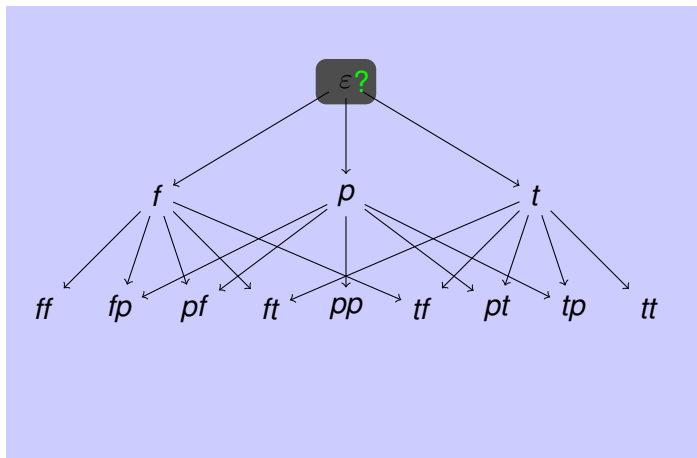
## Discrete Event Systems

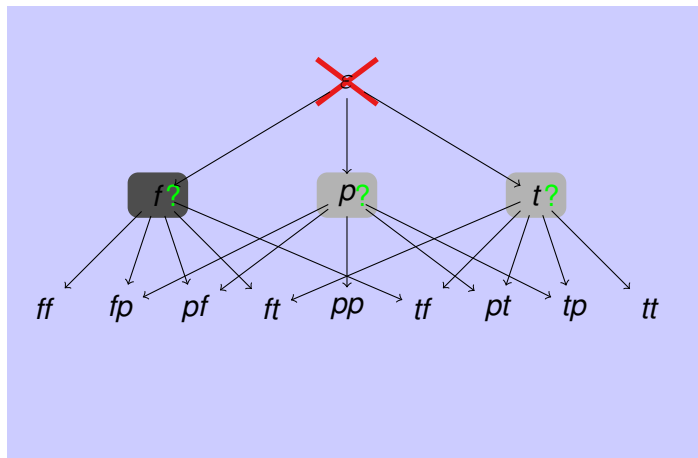
- $\mathcal{L}_h = \{\omega_h\}$  is a finite word defined on  $\Sigma_f$
- $h$  is a candidate iff

$$\mathcal{L}_M \cap \mathcal{L}_O \cap \mathcal{L}_h \neq \emptyset$$

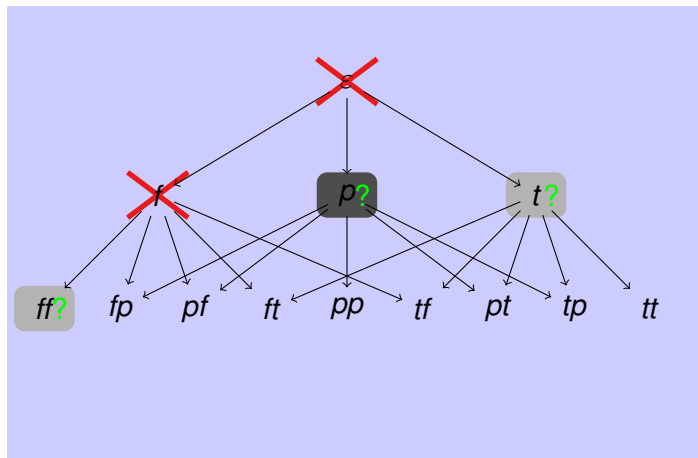




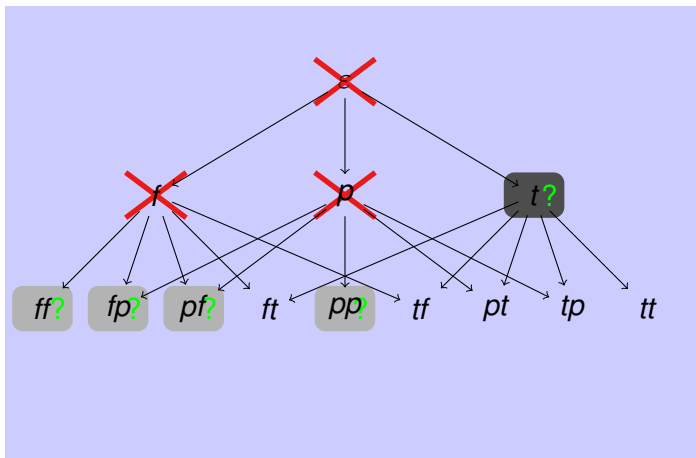


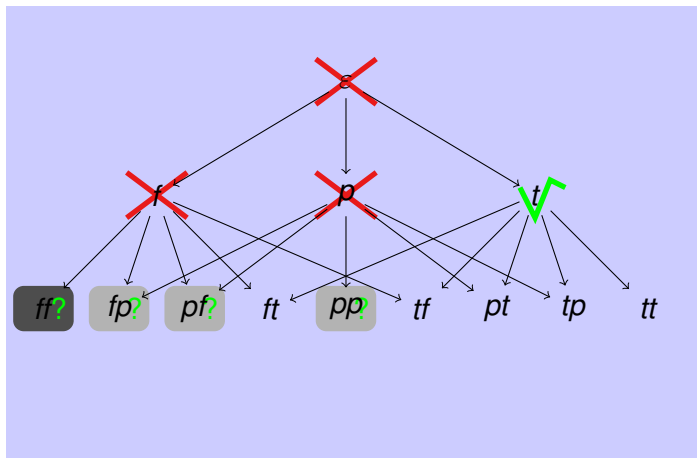


Successors of hypothesis  $h$  is all its children



But ignore successors that are covered by existing hypotheses





Also: termination issue (not discussed here)

## Principle

- If hypothesis  $h$  is not a candidate, the output is not very informative

A conflict is a generalisation of a test failure:

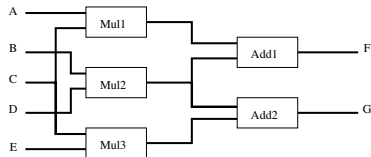
- Why did the test fail?

How to use conflicts:

- An earlier conflict may discard a new hypothesis
- Conflicts can reduce the set of successors

# Conflict Example

## Static System



$$\begin{array}{ll} A = B = E = 3 & F = 10 \\ C = D = 2 & G = 12 \end{array}$$

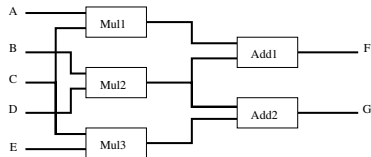
Testing if no component is abnormal:

$$\begin{array}{l} \Phi_M, \Phi_O, \\ (\neg Ab(Mul1) \wedge \neg Ab(Mul2) \wedge \neg Ab(Mul3) \\ \wedge \neg Ab(Add1) \wedge \neg Ab(Add2)) \end{array} \stackrel{?}{\models} \perp$$



# Conflict Example

## Static System



$$A = B = E = 3$$

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$$C = D = 2$$

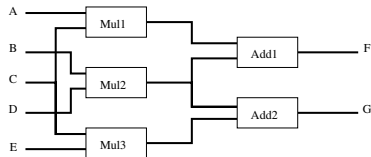
$$G = 12$$

Testing if no component is abnormal:

$$\begin{aligned} & \Phi_M, \Phi_O, \\ & \neg Ab(Mul1), \neg Ab(Mul2), \neg Ab(Mul3), \quad ? \\ & \neg Ab(Add1), \neg Ab(Add2) \quad \models \perp \end{aligned}$$

# Conflict Example

## Static System



$$A = B = E = 3$$

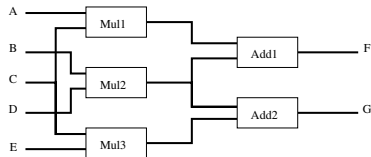
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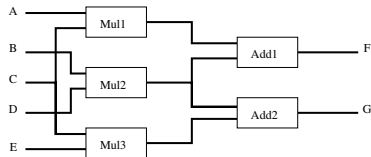
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Testing if no component is abnormal:

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Three successors:

- Only component *Mul1* is abnormal
- Only component *Mul2* is abnormal
- Only component *Add1* is abnormal



$$A = B = E = 3$$

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Three successors:

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If hypothesis  $h$  is not a candidate, then

$$\mathcal{L}_M \cap \mathcal{L}_O \cap \mathcal{L}_h = \emptyset \quad (1)$$

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We reformulate  $\mathcal{L}_h = \mathcal{L}_0 \cap \dots \cap \mathcal{L}_k$

$$\mathcal{L}_M \cap \mathcal{L}_O \cap \mathcal{L}_0 \cap \dots \cap \mathcal{L}_k = \emptyset \quad (2)$$

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For some  $C = \{C_0, \dots, C_{k'}\} \subseteq \{0, \dots, k\}$  (we prefer  $C$  as small as possible),

$$\mathcal{L}_M \cap \mathcal{L}_O \cap \mathcal{L}_{C_0} \cap \dots \cap \mathcal{L}_{C_{k'}} = \emptyset$$

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**$C = \text{conflicts}$**



# Example



## Discrete Event System

$$\Sigma_f = \{a, b, c\} \text{ and } \mathcal{L}_h = \{a\}$$

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$$\{a\} = \mathcal{L}_0 \cap \mathcal{L}_1 \cap \mathcal{L}_2 \cap \mathcal{L}_3 \cap \mathcal{L}_4 \cap \mathcal{L}_5$$

- $\mathcal{L}_0 = \Sigma_f^* a \Sigma_f^*$
- $\mathcal{L}_1 = (\Sigma_f^*) \setminus (\Sigma_f^* a \Sigma_f^* a \Sigma_f^*)$
- $\mathcal{L}_2 = (\Sigma_f^*) \setminus (\Sigma_f^* a \Sigma_f^* b \Sigma_f^*)$
- $\mathcal{L}_3 = (\Sigma_f^*) \setminus (\Sigma_f^* a \Sigma_f^* c \Sigma_f^*)$
- $\mathcal{L}_4 = (\Sigma_f^*) \setminus (\Sigma_f^* b \Sigma_f^* a \Sigma_f^*)$
- $\mathcal{L}_5 = (\Sigma_f^*) \setminus (\Sigma_f^* c \Sigma_f^* a \Sigma_f^*)$

$$\Sigma_f = \{a, b, c\} \text{ and } \mathcal{L}_h = \{a\}$$

$$\{a\} = \mathcal{L}_0 \cap \mathcal{L}_1 \cap \mathcal{L}_2 \cap \mathcal{L}_3 \cap \mathcal{L}_4 \cap \mathcal{L}_5$$

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- $\mathcal{L}_1 = (\Sigma_f^*) \setminus (\Sigma_f^* a \Sigma_f^* a \Sigma_f^*)$
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- $\mathcal{L}_4 = (\Sigma_f^*) \setminus (\Sigma_f^* b \Sigma_f^* a \Sigma_f^*)$

Conflict:  $\{\mathcal{L}_0, \mathcal{L}_1, \mathcal{L}_3, \mathcal{L}_4\}$

Successors:  $aa$ ,  $ac$ , and  $ba$

# Example

## Discrete Event System

$$\Sigma_f = \{a, b, c\} \text{ and } \mathcal{L}_h = \{a\}$$

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- $\mathcal{L}_0 = \Sigma_f^* a \Sigma_f^*$
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Conflict:  $\{\mathcal{L}_0, \mathcal{L}_1, \mathcal{L}_3, \mathcal{L}_4\}$

Successors:  $aa$ ,  $ac$ , and  $ba$

$\Sigma_f = \{a, b, c\}$  and  $\mathcal{L}_h = \{ab\}$

Conflict:

- $\mathcal{L}_i = (\Sigma_f^*) \setminus (\Sigma_f^* b \Sigma_f^* b \Sigma_f^*)$
- $\mathcal{L}_j = (\Sigma_f^*) \setminus (\Sigma_f^* c \Sigma_f^*)$

Successors: *abb, bab, abc, acb, and cab*

$\Sigma_f = \{a, b, c\}$  and  $\mathcal{L}_h = \{ab\}$

Conflict:

- $\mathcal{L}_i = (\Sigma_f^*) \setminus (\Sigma_f^* b \Sigma_f^* b \Sigma_f^*)$
- $\mathcal{L}_j = (\Sigma_f^*) \setminus (\Sigma_f^* c \Sigma_f^*)$

Successors:  $abb$ ,  $bab$ ,  $abc$ ,  $acb$ , and  $cab$

- Given a hypothesis  $h$ , define properties
  - $p_{\text{desc}}(h)$ : property satisfied by all hypotheses  $h' \succeq h$
  - $p_{\text{dese}}(h)$ : property satisfied by all hypotheses  $h' \not\preceq h$
  
- A possible decomposition of  $\{h\}$ :
  - $p_{\text{desc}}(h)$
  - $\forall h' \in \text{children}(h), p_{\text{dese}}(h')$
  
- $C = \{p_1, \dots, p_k\}$  is a conflict for  $h$  iff
  - $\forall h' : p_{\text{desc}}(h') \in C \Rightarrow h' \preceq h$
  - $\forall h' : p_{\text{dese}}(h') \in C \Rightarrow h' \not\preceq h$
  
- Successors of conflict  $C = \{p_1, \dots, p_k\}$ 
  - Let  $\Omega = \{h' \mid p_{\text{dese}}(h') \in C\}$
  - Successors:  $\bigcup_{h' \in \Omega} (h \otimes h')$

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## Diagnosis Problem

- Electricity transmission network
- Alarm log
- Hypothesis: a sequence of “unexplained” events

## Metrics

- Number of components: 3 to 105
- Component model:
  - 8 to 1,024 (more often) states
  - 44 to 92,800 transitions
- Number of minimal candidates: up to 27 and more

	<i>N</i>	<i>M</i>	<i>C</i>	<i>A</i>	PF	JT
window-250	1	0	2	3	0.3	1.5
chunk-004	1	2	3	3	0.8	2
chunk-056	1	4	4	7	1.7	2.6
window-618	1	0	6	2	0.7	–time–
window-527	2	1	11	8	2.7	–time–
window-347	4	9	32	13	106.1	–time–
window-336	?	?	58	49	–time–	–time–
window-335	?	?	67	66	–time–	–time–
chunk-089	?	?	105	146	–time–	–memory–
window-410	?	?	19	13	–time–	5
window-409	?	?	22	14	–time–	5.3
Nb problems solved (/129)					<b>116</b>	<b>35</b>

*N*: number of minimal candidates,

*M*: maximum number of faults in a minimal candidate,

*C*: number of components in the problem,

*A*: number of alarms,

PF: runtime for PF running SAT, and

JT: runtime for automata-based approach (in seconds)

## Contribution

A generalised perspective of conflicts for non trivial hypothesis search space.

## Extensions

- Application to hybrid systems
- Conflicts = explanations