

DATA 61

Solving Diagnosability of Hybrid Systems
via Abstraction and Discrete-Event Techniques:

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Diagnosability of Hybrid Systems

Problem Definition



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A **hybrid system** is a system that involves both continuous (state) and discrete (mode) dynamics.

We assume a strong-fault model (some knowledge on the faulty behaviour).

A system is **diagnosable** if the occurrence of every fault can always be detected and identified by an observer.



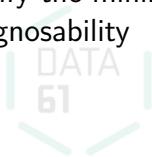
Our Contributions



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Diagnosability of HS with DES Techniques:

- We discretise the system.
- We use DES techniques to prove diagnosability.
- We use an incremental approach to identify the minimal amount of information. necessary for diagnosability



Outline

Running Example

Discretisation of Hybrid Systems

The General Idea

Discernibility

Ephemerality

Diagnosability of DES

Incremental Approach

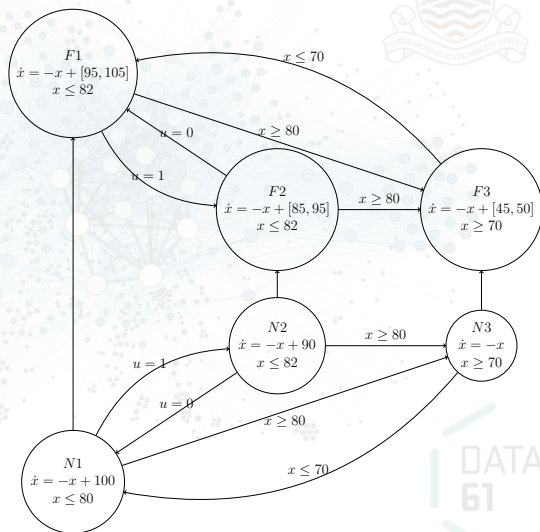
Concluding Remarks



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An Example



Observations: $y = x$, $\dot{y} = \dot{x}$

Outline



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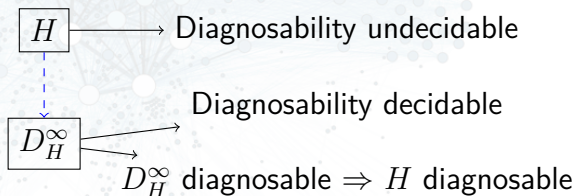
Concluding Remarks



General Idea



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Discretisation from H to D_H^∞



How to compute D_H^∞ :

- Keep the set of modes and transitions (including loops on every mode)
- Compute **discernibility** between modes (discretise the observation)
- Compute **ephemerality** of sets of modes



Discernibility

Definition

Two modes m_1 and m_2 are **discernible** if the observations always allow to determine that you are not in m_2 when you are in m_1 (and vice-versa)



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Indicators

Definition



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An **indicator** is a constraint on the observable variables

Three possible interactions between a mode and an indicator:

- the indicator **always** holds in the mode
- the indicator **never** holds in the mode
- the indicator **sometimes** holds in the mode



Discernibility and Indicators



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If an indicator

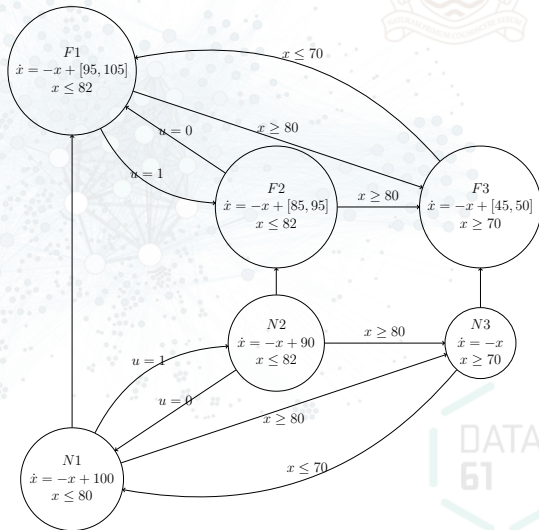
- always holds in m_1 and
- never holds in m_2

then m_1 and m_2 are discernible.



Example

Discernibility (1/4)



Observations: $y = x, \dot{y} = \dot{x}$

Example

Discernibility (2/4)



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- $indi_1: \dot{y} + y = 100$ (derived from mode $N1$)
- $indi_2: \dot{y} + y = 90$ (mode $N2$)
- $indi_3: \dot{y} + y = 0$ (mode $N3$)
- $indi_4: \dot{y} + y \in [95, 105]$ (mode $F1$)
- $indi_5: \dot{y} + y \in [85, 89]$ (mode $F2$)
- $indi_6: \dot{y} + y \in [45, 50]$ (mode $F3$)



Example

Discernibility (3/4)



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Indicator function:

1 : the indicator is always satisfied in this mode

-1 : the indicator is never satisfied in this mode

0 : the indicator is sometimes satisfied in this mode

	<i>N1</i>	<i>N2</i>	<i>N3</i>	<i>F1</i>	<i>F2</i>	<i>F3</i>
<i>indi</i> ₁	1			0		
<i>indi</i> ₂		1				
<i>indi</i> ₃			1			
<i>indi</i> ₄	1			1		
<i>indi</i> ₅					1	
<i>indi</i> ₆						1

(-1s not represented)



Example

Discernibility (4/4)

Indiscernibility matrix:

- two modes are discernible if

$$\{L(m_1, indi), L(m_2, indi)\} = \{-1, 1\}$$

for some indicator *indi*

	N1	N2	N3	F1	F2	F3
N1	1			1		
N2		1				
N3			1			
F1	1			1		
F2					1	
F3						1

a 1 indicates that the modes are not discernible



Diagnosability of DES



Twin Plant:

- Make a copy D' of D
- Remove the faulty modes of D'
- Synchronise D with D' (mode-based observations: **remove discernible pairs**)

A **counter-example** is a cycle in the twin plant that is

- reachable,
- and ambiguous.

Theorem: If there is no counter-example the DES is diagnosable

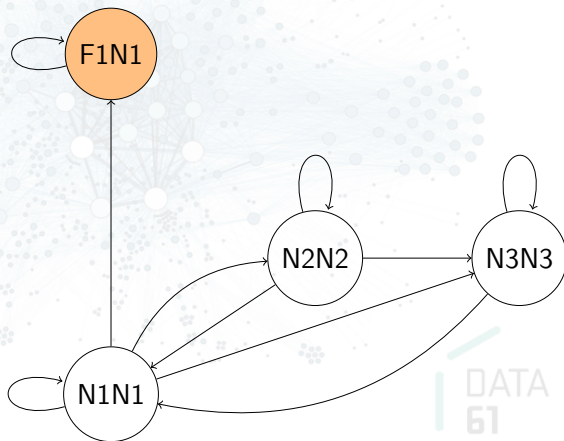


Twin Plant for D_H^∞

Running Example



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Issue:

- The discretisation inserts loops on every mode: how do we model the fact that the system cannot stay in this mode?

Definition:

- A set of modes is **ephemeral** if the system cannot stay forever in this set of modes



Ephemerality

Running Example

- In mode $N1$,
 - Derivative: $\dot{x} = 100 - x \geq 20$
 - Invariant: $x \leq 80$ \Rightarrow eventually the system must leave mode $N1$
 $\{N1\}$ is ephemal



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Ephemerality

Running Example



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- In mode $N1$,
 - Derivative: $\dot{x} = 100 - x \geq 20$
 - Invariant: $x \leq 80$
- \Rightarrow eventually the system must leave mode $N1$
- $\{N1\}$ is ephemal

The ephemeral sets include:

- $\{N1, N2\}$
- $\{N3\}$
- $\{F1, F2\}$
- $\{F3\}$.



Diagnosability of DES D

Theory



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Twin Plant:

- Make a copy D' of D
- Remove the faulty modes of D'
- Synchronise D with D' (mode-based observations: **remove discernible pairs**)

A **counter-example** is a cycle in the twin plant that is

- reachable,
- fair (non-ephemeral),
- and ambiguous.

Theorem: The DES is diagnosable **iff** there is no counter-example

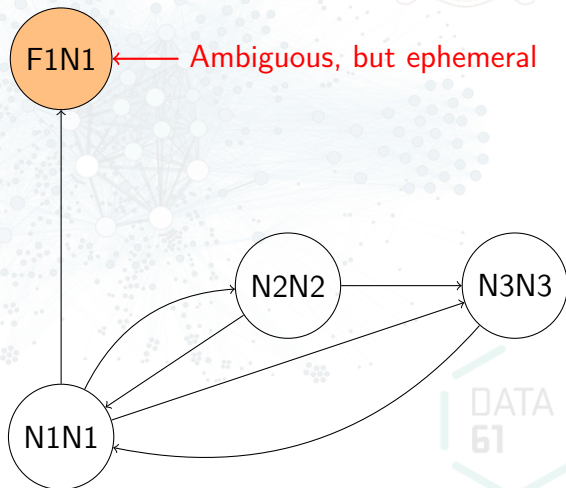


Twin Plant for D_H^∞

Running Example



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Motivation

Why Use Incremental Approach?



- Computing all indicators is expensive
 - They are exponentially many
 - Many indicators are useless/redundant
- Computing the ephemeral sets is expensive
- Using all indicators during diagnosis is expensive
 - We want to identify the indicators that are helpful



General Idea

H → Diagnosability undecidable

D_H^∞ → Diagnosability decidable

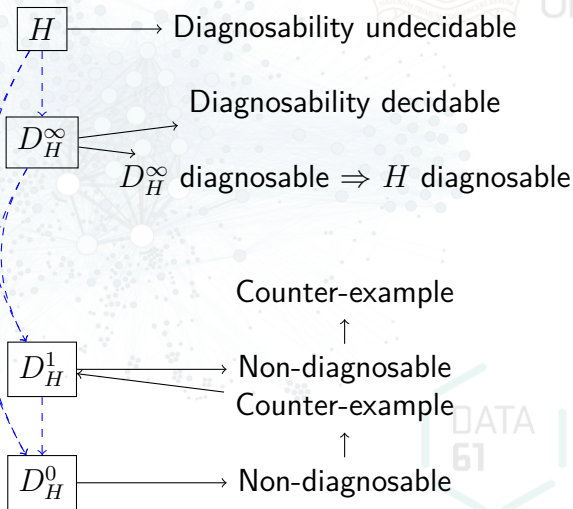
D_H^∞ diagnosable $\Rightarrow H$ diagnosable



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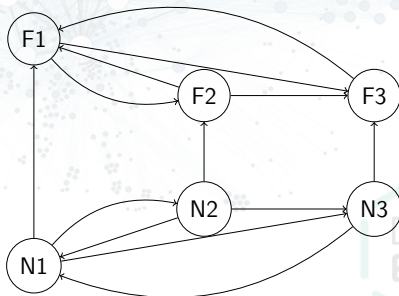
General Idea



0. Maximal Abstraction: D_H^0

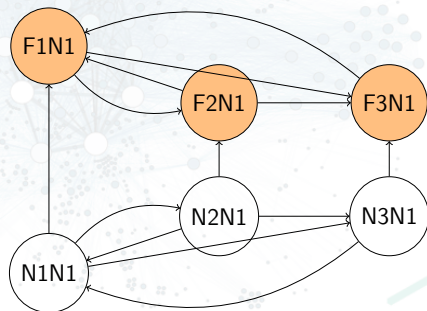
How to Compute D_H^0 :

- Keep the modes and transitions (including loops)
- Ignore state dynamics and guards
- No indiscernibility, no ephemerality



0. Twin Plant

Running Example



0. Counter Example

Running Example



If the system generates the following faulty behaviour:

$$b_F = N1 \rightarrow F1(\rightarrow F1)^\infty$$

then the diagnoser might believe that what is happening is:

$$b_N = N1 \rightarrow N1(\rightarrow N1)^\infty$$



0. Negating the Counter Example



Critical Pair (reminder):

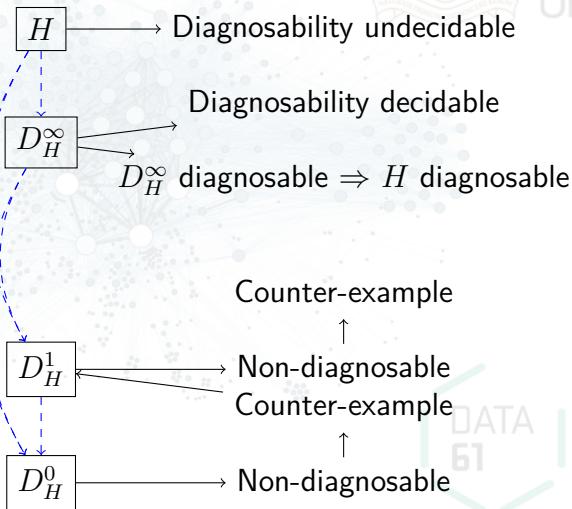
- $b_F = N1 \rightarrow F1(\rightarrow F1)^\infty$
- $b_N = N1 \rightarrow N1(\rightarrow N1)^\infty$

Answer:

- $\{F1\}$ is ephemeral, therefore the counter example is not valid.



Back to the General Idea

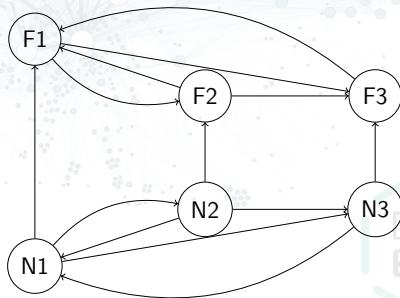


1. New Abstraction: D_H^1

Running Example

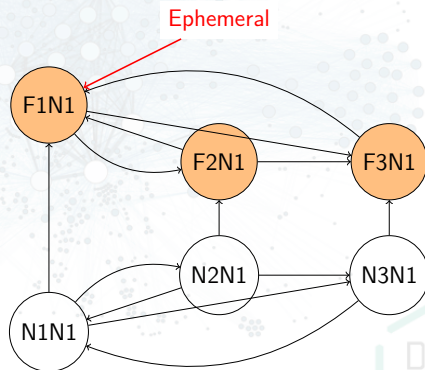
Known ephemeral sets

- $\{F1\}$



1. Twin Plant

Running Example



1. Counter Example

Running Example



If the system generates the following faulty behaviour:

$$b_F = N1 \rightarrow F1 \rightarrow F2(\rightarrow F1 \rightarrow F2)^\infty$$

then the diagnoser might believe that what is happening is:

$$b_N = N1 \rightarrow N1 \rightarrow N1(\rightarrow N1 \rightarrow N1)^\infty$$

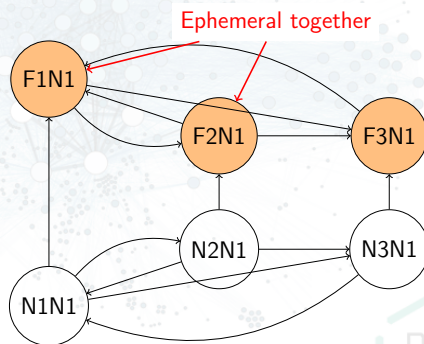
Answer:

- It can be shown that $\{F1, F2\}$ is ephemeral.



2. Twin Plant

Running Example



2. Discernibility

Running Example



If the system generates the following faulty behaviour:

$$b_F = N1 \rightarrow F1 \rightarrow F2 \rightarrow F3(\rightarrow F1 \rightarrow F2 \rightarrow F3)^\infty$$

then the diagnoser might believe that what is happening is:

$$b_N = N1 \rightarrow N1 \rightarrow N1 \rightarrow N1(\rightarrow N1 \rightarrow N1 \rightarrow N1)^\infty$$

Answer:

- Indicator $I1$ always discern $N1$ from $F2$

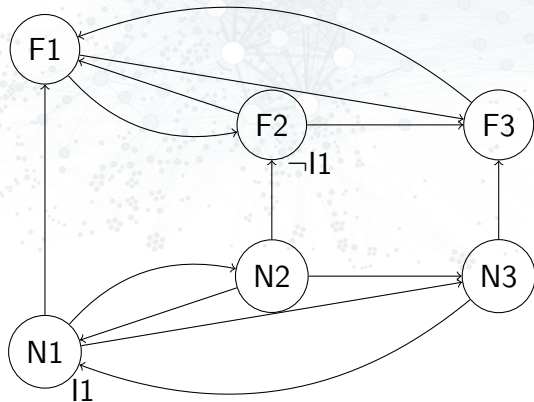


3. New Abstraction: D_H^3

Running Example

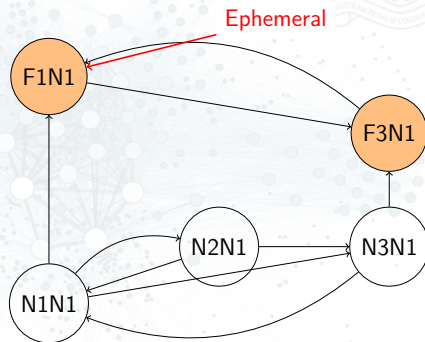
Known ephemeral sets

- $\{F1, F2\}$



3. New Twin Plant: D_H^3

Running Example



New counter example:

- $b_F = N1 \rightarrow F1 \rightarrow F2 \rightarrow F3 (\rightarrow F1 \rightarrow F2 \rightarrow F3)^\infty$
- $b_N = N1 \rightarrow N1 \rightarrow N2 \rightarrow N1 (\rightarrow N1 \rightarrow N2 \rightarrow N1)^\infty$

etc.

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Summarising the Approach



- To check diagnosability of hybrid systems, we discretise the hybrid model:
 - we keep the list of modes
 - we keep the list of transitions
 - we infer *ephemerality* properties
 - we infer *discernibility* properties between modes
- We compute a subset of these properties sufficient for diagnosability → near-optimal observability





- Ephemerality and discernibility.
How to compute these properties ?
- “ D_H^∞ not diagnosable” does not imply “ H not diagnosable”.
What can we do if D_H^∞ is not diagnosable ?
- Symbolic tools.
Using BDDs to verify diagnosability of networks of systems
with $> 2^{100}$ modes.

